

ON THE COMPUTABILITY-THEORETIC COMPLEXITY OF TRIVIAL, STRONGLY MINIMAL MODELS

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ABSTRACT. We show the existence of a trivial, strongly minimal (and thus uncountably categorical) theory for which the prime model is computable and each of the other countable models computes $\mathbf{0}''$. This result shows that the result of Goncharov/Harizanov/Laskowski/Lempp/McCoy [GHLLM03] is best possible for trivial strongly minimal theories in terms of computable model theory. We conclude with some remarks about axiomatizability.

Both vector spaces and algebraically closed fields have the property that the uncountable structures in these classes are determined up to isomorphism by their cardinality. Model theory provides a general framework in which to study this behavior. A first order theory T is called κ -categorical (for an infinite cardinal κ) if $\mathcal{A} \cong \mathcal{B}$ whenever $\mathcal{A}, \mathcal{B} \models T$ and $|\mathcal{A}| = |\mathcal{B}| = \kappa$. Morley [Mo65] proved that if T is categorical in some uncountable cardinality, then it is categorical in all uncountable cardinalities. (We assume here and for the rest of this paper that the language is countable.) Therefore, any ω_1 -categorical theory has a unique model of each uncountable cardinality.

Baldwin and Lachlan [BL71] showed that for an uncountably categorical theory T , either there is a unique countable model of T , or the countable models of T form an elementary chain of length $\omega + 1$: $\mathcal{A}_0 \prec \mathcal{A}_1 \prec \cdots \prec \mathcal{A}_\omega$, where \mathcal{A}_0 is the prime (or atomic) model and \mathcal{A}_ω is the countably saturated model. The examples of vector spaces (over a fixed countable field) and algebraically closed fields fall into this second category, with the models in the elementary chain determined by the cardinality of a basis and by the transcendence degree. It is this case of uncountably categorical but not ω -categorical theories that we are interested in.

Together, Morley's Categoricity Theorem and the Lachlan-Baldwin Theorem give a good algebraic picture of the infinite models of an uncountably categorical theory T . However, there are additional questions to consider about the complexity of such theories and models. For example, from a model-theoretic point of view, one can ask about the complexity of sentences required to axiomatize these theories,

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and from a computability-theoretic point of view, one can ask about the complexity of constructing the models of such a theory.

To address the latter question, we recall that a countable model is said to be *decidable* if it has an isomorphic copy with universe ω for which the elementary diagram forms a computable set of formulas; and *computable* if the open diagram of some such copy forms a computable set of formulas. (We assume that we are working in a *computable* language, i.e., that the arity of each relation and function symbol in the language is uniformly computable. Note that these notions easily relativize to the notions of “ \mathbf{d} -decidable” and “ \mathbf{d} -computable” for any Turing degree \mathbf{d} .)

Harrington [Ha74] and Khisamiev [Kh74] showed that all countable models of a decidable uncountably categorical theory are decidable. It follows trivially from the definitions that if an uncountably categorical theory has a decidable model then the theory is decidable, and so all its countable models are decidable.

In contrast, Goncharov [Go78] first exhibited an uncountably categorical theory for which one but not all countable models are computable; for his theory, the prime model is computable while all other countable models are only $\mathbf{0}'$ -computable. More examples of uncountably categorical theories for which some but not all countable models are computable were found by Kudaibergenov [Ku80], Khoussainov/Nies/Shore [KNS97], Nies [Ni99] and others. These theories have three properties in common: They use infinite languages, they are trivial and strongly minimal (defined below), and all the countable models are computable from $\mathbf{0}'$.

Herwig/Lempp/Ziegler [HLZ99] showed that the use of an infinite language is not essential by constructing a theory in a language of three binary relation symbols for which the prime model is computable but all nonprime models are only $\mathbf{0}'$ -computable. Moreover, they showed that for any Turing degree $\mathbf{a} \leq_T \mathbf{0}'$, there is a trivial strongly minimal theory T in this language such that the prime model of T is computable while (the open diagrams of) all the nonprime models have copies in exactly the degrees $\geq_T \mathbf{a}$, i.e., the nonprime models all “code” the degree \mathbf{a} .

The second common property of these theories is that each example is trivial and strongly minimal. A theory T is called *strongly minimal* if every definable subset (with parameters) of each model $\mathcal{A} \models T$ is either finite or cofinite. Such a model $\mathcal{A} \models T$ is *trivial* (or *has trivial pregeometry*) if the algebraic closure of any subset $D \subseteq A$, denoted by $\text{acl}(D)$, is the union of the algebraic closures of the elements of D . (This property depends only on the theory of \mathcal{A} . An element $a \in A$ is *algebraic over* D if there are parameters $\bar{d} \in D$ and a formula $\psi(x, \bar{d})$ such that $\mathcal{A} \models \psi(a, \bar{d})$ and the set of all $b \in A$ for which $\mathcal{A} \models \psi(b, \bar{d})$ is finite. The *algebraic closure* of D is the set of all elements which are algebraic over D . In the case of algebraically closed fields and vector spaces, $\text{acl}(D)$ is the algebraic closure of the subfield generated by D , and the span of D , respectively. Buechler [Bu96] is a good reference on strongly minimal theories.) Strong minimality implies uncountable categoricity, so this property shows that it is possible to get the behaviors described above within a restricted class of uncountably categorical theories.

The third property is our main concern in this paper. Although all of the models in the examples above are $\mathbf{0}'$ -computable, it was not clear for several years why there should be any fixed arithmetic bound on the complexity of these models. The existence of such an upper bound was given by a purely model-theoretic result by Goncharov/Harizanov/Laskowski/Lempp/McCoy [GHLLM03]: If T is a trivial

strongly minimal theory, $\mathcal{M} \models T$ and $\text{Th}(\mathcal{M}_M)$ denotes the theory of \mathcal{M} in the expanded language which includes a constant symbol for each element of \mathcal{M} , then $\text{Th}(\mathcal{M}_M)$ is model complete, and hence $\forall\exists$ -axiomatizable (see, e.g., Chang/Keisler [CK90, pp. 186-187]). Therefore, if $\mathcal{M} \models T$ is computable, then both $\text{Th}(\mathcal{M}_M)$ and $\text{Th}(\mathcal{M}) = T$ are $\mathbf{0}''$ -decidable. By the relativized version of the theorem of Harrington [Ha74] and Khisamiev [Kh74], all countable models of T are $\mathbf{0}''$ -decidable. (This statement is stronger than saying they are $\mathbf{0}''$ -computable; it says that each model has an isomorphic copy in which the elementary diagram (and not just the open diagram) is $\mathbf{0}''$ -computable.)

There is a computability-theoretic gap between the results of Goncharov [Go78] (in which the countable nonprime models of T are computable in $\mathbf{0}'$) and of Goncharov/Harizanov/Laskowski/Lempp/McCoy [GHLLM03] (in which the upper bound of $\mathbf{0}''$ is given). Is it possible for a trivial strongly minimal theory to have a computable model and another countable model which is only $\mathbf{0}''$ -computable? In this paper, we give an affirmative answer and thereby show that the result of Goncharov/Harizanov/Laskowski/Lempp/McCoy [GHLLM03] is best possible in terms of computable model theory (i.e., the complexity of the countable models).

Theorem 1. *There is a trivial, strongly minimal (and thus uncountably categorical) first-order theory T such that the prime model of T is computable and $\mathbf{0}''$ is computable from the atomic diagram of any nonprime countable model of T .*

We prove this theorem in section 1. As a corollary to this theorem, we can also conclude that the $\forall\exists$ -axiomatization of the theory of a strongly minimal model in the expanded language naming constants is optimal, i.e., there is no axiomatization by Boolean combinations of \forall -formulas. However, this can be seen already from much simpler examples, e.g., the theory of an equivalence relation with all equivalence classes of size two. Section 2 contains some general remarks about axiomatizability by $\forall\exists$ -formulas, \forall -formulas, and Boolean combinations of \forall -formulas.

1. THE PROOF OF THEOREM 1

We begin by defining the type of family of sets which we will code into our theory.

Definition 2. An infinite family \mathcal{F} of finite subsets of ω is an **almost everywhere** (or **a.e.**) **family** if for every $n \in \omega$, either $n \in X$ for all but finitely many $X \in \mathcal{F}$ or $n \notin X$ for all but finitely many $X \in \mathcal{F}$. Such a family \mathcal{F} has an **infinite limit** if the set

$$\begin{aligned} I_{\mathcal{F}} &= \{n \in \omega \mid n \in X \text{ for all but finitely many } X \in \mathcal{F}\} \\ &= \{n \in \omega \mid n \in X \text{ for infinitely many } X \in \mathcal{F}\} \end{aligned}$$

is infinite.

To code this type of family, we generalize the notion of ‘‘cubes’’ introduced by Khossainov/Nies/Shore [KNS97]. Fix a language \mathcal{L} consisting of binary relation symbols F_i (for $i \in \omega$), which we will all assume to be symmetric and irreflexive and to have in-degree and out-degree ≤ 1 , so we can think of them as coding permutations f_i of order 2 on subsets of the model. (If we used function symbols f_i in our language rather than relation symbols, then by the comments in Section 2, the theory $\text{Th}(\mathcal{M}_M)$ for a model $\mathcal{M} \models T_{\mathcal{F}}$ (defined below) would have a quantifier-depth 1 axiomatization and hence the most we can code into such a model is $\mathbf{0}'$.)

Before giving the formal definition of an S -cube (for $S \subseteq \omega$), we give several concrete examples. Suppose $S = \{n_0\}$. In this case, an S -cube consists of two elements a_0 and b_0 such that $F_{n_0}(a_0, b_0)$ (and hence $F_{n_0}(b_0, a_0)$ by symmetry) and $\neg F_m(a_0, b_0)$ for all $m \neq n_0$ (and hence $\neg F_m(b_0, a_0)$ by symmetry). Notice that the function f_{n_0} induced by F_{n_0} and defined by $f_{n_0}(a_0) = b_0$ and $f_{n_0}(b_0) = a_0$ gives an \mathcal{L} -automorphism of the S -cube.

Next, suppose $S = \{n_0, n_1\}$. To construct the S -cube, we begin with two $\{n_0\}$ -cubes consisting of the elements a_0 and b_0 connected by $F_{n_0}(a_0, b_0)$ and the elements a_1 and b_1 connected by $F_{n_0}(a_1, b_1)$. We join these two $\{n_0\}$ -cubes by adding the relations $F_{n_1}(a_0, a_1)$ and $F_{n_1}(b_0, b_1)$ and no other relations (except those forced by symmetry). Notice that the function f_{n_1} induced by F_{n_1} and defined by $f_{n_1}(a_i) = a_{1-i}$ and $f_{n_1}(b_i) = b_{1-i}$ for $i \leq 1$ is both an \mathcal{L} -isomorphism between the two $\{n_0\}$ -cubes and an \mathcal{L} -automorphism of the S -cube. Similarly, the function f_{n_0} induced by F_{n_0} is an \mathcal{L} -automorphism of the S -cube. Combining these automorphisms, it is clear that any two elements of the S -cube are automorphic. (We could get an isomorphic S -cube by starting with two $\{n_1\}$ -cubes and connecting them using F_{n_0} . Therefore, the roles played by n_0 and n_1 are actually symmetric.)

For one last example, suppose $S = \{n_0, n_1, n_2\}$. To construct the S -cube, we begin with two $\{n_0, n_1\}$ -cubes given by a_0, a_1, b_0 and b_1 (with relations as above) and by a'_0, a'_1, b'_0 and b'_1 (with relations defined similarly). We join these two cubes by adding the relations $F_{n_2}(a_i, a'_i)$ and $F_{n_2}(b_i, b'_i)$ for $i \leq 1$ and no other relations (except those forced by symmetry). As above, the function f_{n_2} induced by F_{n_2} is both an \mathcal{L} -isomorphism between the $\{n_0, n_1\}$ -cubes and an \mathcal{L} -automorphism of the S -cube. Similarly, the functions f_{n_0} and f_{n_1} induced by F_{n_0} and F_{n_1} are \mathcal{L} -automorphisms of the S -cube. Combining these automorphisms, it is clear that any two elements of the S -cube are automorphic.

We now give the formal definition of an S -cube.

Definition 3. Given a nonempty (finite or infinite) subset $S \subseteq \omega$, we let $S = \{n_0 < n_1 < \dots\}$ be a (not necessarily effective) enumeration of S in order, and set $S_k = \{n_0, \dots, n_k\}$ (for $k < |S|$ if S is finite). We define the notion of an S -cube by induction on k .

An S_0 -cube is a 2-element subset $\{x, y\}$ of an \mathcal{L} -model such that $F_{n_0}(x, y)$ (coding $f_{n_0}(x) = y$ and $f_{n_0}(y) = x$) and $\neg F_n(x, y)$ for all $n \neq n_0$.

An S_{k+1} -cube is the disjoint union of two S_k -cubes C_0 and C_1 such that $F_{n_{k+1}}$ induces an \mathcal{L} -isomorphism between C_0 and C_1 , and such that $\neg F_n(x_0, x_1)$ for all $n \neq n_{k+1}$, all $x_0 \in C_0$ and all $x_1 \in C_1$.

For infinite S , we define an S -cube to be the “limit” of the S_k -cubes, i.e., an infinite subset A of an \mathcal{L} -model such that for all $k \in \omega$, every element $x \in A$ is an element of a unique S_k -cube $C \subset A$ and such that any two elements $x_0, x_1 \in A$ are connected by a finite chain of F_i -relations.

Notice that if $X \subsetneq Y$, then any Y -cube contains X -cubes. We will use the term X -component to refer to an X -cube which is not contained in a larger cube. (Intuitively, an X -component is a connected component of the model viewed as a graph.) There are two important properties of an S -component. First, any two elements of an X -component are automorphic. Second, if x is an element of an S -component, then there exists a y such that $F_n(x, y)$ if and only if $n \in S$.

For the rest of this paper, \mathcal{F} denotes an a.e. family of finite subsets of ω with an infinite limit $I_{\mathcal{F}}$. We code such a family into an \mathcal{L} -theory as follows. Let $\mathcal{A}_{0, \mathcal{F}}$

denote the \mathcal{L} -structure consisting of a single X -component for each $X \in \mathcal{F}$ and let $T_{\mathcal{F}}$ be the theory of this structure. We first examine some model theoretic properties of such a theory.

Lemma 4. *$\mathcal{A}_{0,\mathcal{F}}$ is the atomic model of $T_{\mathcal{F}}$. Furthermore, for any nonatomic model $\mathcal{M} \models T_{\mathcal{F}}$ and any $m \in \mathcal{M} \setminus \mathcal{A}_{0,\mathcal{F}}$, m sits inside an $I_{\mathcal{F}}$ -component. (Thus any model of $T_{\mathcal{F}}$ can be denoted as $\mathcal{A}_{\kappa,\mathcal{F}}$ and is the disjoint union of the atomic model and κ many $I_{\mathcal{F}}$ -components.)*

Proof. To show that $\mathcal{A}_{0,\mathcal{F}}$ is the atomic model of $T_{\mathcal{F}}$, fix $a \in \mathcal{A}_{0,\mathcal{F}}$ and $X \in \mathcal{F}$ such that a is an element of the X -component in $\mathcal{A}_{0,\mathcal{F}}$. Since $I_{\mathcal{F}}$ is infinite, we can fix $i \in I_{\mathcal{F}} \setminus X$. Since \mathcal{F} is an a.e. family, there are finitely many $Y \in \mathcal{F}$ such that $i \notin Y$, and we let X_0, X_1, \dots, X_{k-1} list these elements of \mathcal{F} which are not equal to X . For each $0 \leq j < k$, fix n_j such that either $n_j \in X \setminus X_j$ or $n_j \in X_j \setminus X$. Let $\psi_a(x)$ be the formula that says $\exists y(F_m(x, y))$ for each $m \in X$, $\neg \exists y(F_i(x, y))$ for our fixed element $i \in I_{\mathcal{F}} \setminus X$, and $\neg \exists y(F_{n_j}(x, y))$ for each $0 \leq j < k$ such that the fixed element n_j satisfies $n_j \in X_j \setminus X$.

Clearly, a satisfies $\psi_a(x)$. Moreover, if $b \in \mathcal{A}_{0,\mathcal{F}}$ and b is not in the X -component, then we claim that b does not satisfy $\psi_a(x)$. Fix $Y \in \mathcal{F}$ such that b is in the Y -component of $\mathcal{A}_{0,\mathcal{F}}$. If $i \in Y$, then b fails to satisfy $\neg \exists y(F_i(x, y))$. Suppose $Y = X_j$ for some $0 \leq j < k$. If $n_j \in X \setminus X_j$, then b fails to satisfy the $\exists y(F_{n_j}(x, y))$ conjunct in $\psi_a(x)$. If $n_j \in X_j \setminus X$, then b fails to satisfy the $\neg \exists y(F_{n_j}(x, y))$ conjunct in $\psi_a(x)$. Therefore, the only elements of $\mathcal{A}_{0,\mathcal{F}}$ which satisfy $\psi_a(x)$ are the elements of the X -component.

Thus, for each $k \notin X$, $T_{\mathcal{F}}$ contains the sentence $\neg \exists x, y(\psi_a(x) \wedge F_k(x, y))$. Hence, if \mathcal{M} is any model of $T_{\mathcal{F}}$ and $b \in \mathcal{M}$ satisfies $\psi_a(x)$, then b is an element of an X -component. Since \mathcal{M} must contain such an element b , we have that \mathcal{M} contains an X -component. It follows that $\mathcal{A}_{0,\mathcal{F}}$ embeds into \mathcal{M} , so $\mathcal{A}_{0,\mathcal{F}}$ is the prime (and hence atomic) model of $T_{\mathcal{F}}$. Furthermore, because $T_{\mathcal{F}}$ says that there are exactly $2^{|X|}$ many elements which satisfy $\psi_a(x)$, the principal type satisfied by an element of an X -component is algebraic over \emptyset .

Next, we show that if $\mathcal{M} \models T_{\mathcal{F}}$ is a nonprime model and $m \in \mathcal{M} \setminus \mathcal{A}_{0,\mathcal{F}}$, then m is contained in an $I_{\mathcal{F}}$ -component. Let $I_{\mathcal{F}} = \{i_0 < i_1 < i_2 < \dots\}$ be a (not necessarily effective) enumeration of $I_{\mathcal{F}}$. For each $k \in \omega$, we let C_k denote the Y -cube where $Y = \{i_0, i_1, \dots, i_k\}$.

We think of \mathcal{M} containing an edge between any elements a and b if $F_i(a, b)$ holds for some i and we use graph-theoretic language. We define a modified notion of ‘‘open balls’’ in \mathcal{M} . For any $a \in \mathcal{M}$ and $k \in \omega$, $B(a, k)$ is the set of all $b \in \mathcal{M}$ such that the distance between a and b as measured by the relations F_0, \dots, F_{i_k} is $\leq k$.

By the properties of \mathcal{F} , for any fixed $k \in \omega$, almost all $a \in \mathcal{M}$ satisfy $B(a, k) \cong C_k$. For each $k \in \omega$, the same finite number of exceptional points a for which $B(a, k) \not\cong C_k$ exist in \mathcal{M} as in $\mathcal{A}_{0,\mathcal{F}}$ and hence these exceptional points all come from the finite components of $\mathcal{A}_{0,\mathcal{F}}$ sitting inside \mathcal{M} . Since any point $a \in \mathcal{M} \setminus \mathcal{A}_{0,\mathcal{F}}$ satisfies $B(a, k) \cong C_k$ for all $k \in \omega$, each such a sits inside an $I_{\mathcal{F}}$ -component. Furthermore, since any two points in an $I_{\mathcal{F}}$ -component are automorphic, any two elements in $\mathcal{M} \setminus \mathcal{A}_{0,\mathcal{F}}$ satisfy the same 1-type. Thus, there is a unique nonprincipal (and nonalgebraic) 1-type in \mathcal{M} for any nonprime model \mathcal{M} of $T_{\mathcal{F}}$. \square

Lemma 5. *$T_{\mathcal{F}}$ is trivial and strongly minimal.*

Proof. To see that $T_{\mathcal{F}}$ is strongly minimal, let $\psi(x)$ be any formula, possibly with parameters \bar{b} from a model $\mathcal{M} \models T_{\mathcal{F}}$. We need to show that either $\{m \in \mathcal{M}' \mid \mathcal{M}' \models \psi(m)\}$ or $\{m \in \mathcal{M}' \mid \mathcal{M}' \models \neg\psi(m)\}$ is finite for any elementary extension \mathcal{M}' of \mathcal{M} . Let $\mathcal{L}_{\bar{b}}$ be the restriction of our language to those edge relations occurring in $\psi(x)$. Then the reduct $\mathcal{M}' \upharpoonright \mathcal{L}_{\bar{b}}$ is the disjoint union of (maximal) X -cubes for various finite sets X containing only indices of edge relations from the finite set of indices I_0 of edge relations mentioned in $\psi(x)$, and all but finitely many of these cubes are X_0 -cubes for $X_0 = I_{\mathcal{F}} \cap I_0$ and do not contain a parameter from \bar{b} . Since the stabilizer of \bar{b} in the automorphism group of $\mathcal{M}' \upharpoonright \mathcal{L}_{\bar{b}}$ is 1-transitive on the union U_0 of all those X_0 -cubes, U_0 is a cofinite subset of \mathcal{M}' which is completely contained in either $\{m \in \mathcal{M}' \mid \mathcal{M}' \models \psi(m)\}$ or $\{m \in \mathcal{M}' \mid \mathcal{M}' \models \neg\psi(m)\}$, establishing our claim of strong minimality.

It remains to show that $T_{\mathcal{F}}$ is trivial, or in other words, that $\text{acl}(A) = \bigcup_{a \in A} \text{acl}(a)$ for all $A \subseteq \mathcal{M}$ with $\mathcal{M} \models T_{\mathcal{F}}$. We have already pointed out that $\text{acl}(\emptyset)$ is equal to the set of all finite X -components in \mathcal{M} . (That is, $\text{acl}(\emptyset) = \mathcal{A}_{0, \mathcal{F}} \prec \mathcal{M}$.) We claim that $\text{acl}(A)$ is equal to the set of all finite X -components in \mathcal{M} together with the set of all $I_{\mathcal{F}}$ -components in \mathcal{M} which intersect A . The fact that $T_{\mathcal{F}}$ is trivial follows immediately from this claim.

To see the claim, notice that if $b \in \mathcal{M}$ is in the same $I_{\mathcal{F}}$ -component as $a \in A$, then there is a path in \mathcal{M} from a to b (using various relations F_i) and that b is the unique point at the end of such a path. Therefore, the set of all $I_{\mathcal{F}}$ -components which intersect A are contained in $\text{acl}(A)$. Furthermore, if b is in an $I_{\mathcal{F}}$ -component which does not intersect A , then for any other c in the same $I_{\mathcal{F}}$ -component as b , there is an automorphism of \mathcal{M} which sends b to c and fixes the finite X -components and the $I_{\mathcal{F}}$ -components which intersect A . Therefore, b and c realize the same type over A , this type is not algebraic over A and neither b nor c is in $\text{acl}(A)$. \square

We next examine some computability-theoretic properties of models of $T_{\mathcal{F}}$.

Definition 6. A **c.e. enumeration** of a countable family \mathcal{G} of subsets of ω is a uniformly c.e. sequence of sets X_i such that $\mathcal{G} = \{X_i \mid i \in \omega\}$. If $i \neq j$ implies $X_i \neq X_j$, then we call this enumeration a **c.e. Friedberg enumeration** of \mathcal{G} . (We note that in Russian terminology, a c.e. enumeration is frequently called a **computable numbering**.)

Lemma 7. *If there is a c.e. Friedberg enumeration of \mathcal{F} , then $T_{\mathcal{F}}$ has a computable prime model.*

Proof. Given a c.e. Friedberg enumeration of \mathcal{F} , we construct a computable prime model of $T_{\mathcal{F}}$ by building the X -component for each $X \in \mathcal{F}$ as it is enumerated. That is, we do not need to specify all of the X -component at once, but can build it in stages as the (finitely many) elements of X appear. \square

The following lemma gives the key property for verifying that theories $T_{\mathcal{F}}$ for appropriate families \mathcal{F} can code additional information into their nonprime models.

Lemma 8. *If \mathcal{A} is a nonprime model of $T_{\mathcal{F}}$, then $I_{\mathcal{F}}$ is c.e. relative to the atomic diagram of \mathcal{A} .*

Proof. Fix an element c in one of the $I_{\mathcal{F}}$ -components in \mathcal{A} . Using the atomic diagram of \mathcal{A} , we enumerate the numbers $n \in \omega$ such that there exists an element x in \mathcal{A} which is related to c by F_n . By the definition of the $I_{\mathcal{F}}$ -component, this procedure enumerates $I_{\mathcal{F}}$. \square

By Lemma 8, we can code $0''$ into the nonprime models of \mathcal{F} by making $0''$ computable from any enumeration of $I_{\mathcal{F}}$. We do this by making $I_{\mathcal{F}}$ code the graph of a function which grows fast enough to compute $0''$.

We develop a suitable computable approximation to $0''$. Fix the canonical uniformly c.e. sequence of all c.e. sets W_n for $n \in \omega$ so that $\{n \mid W_n \text{ is finite}\} \equiv_T 0''$. We let $W_{n,s}$ denote the elements of W_n that have been enumerated by stage s . Let $f(n)$ be the least s such that the finite sets among W_0, \dots, W_n have been completely enumerated by stage s .

Lemma 9. *For any function g which dominates f , $0'' \leq_T g \oplus 0'$.*

Proof. Let $k \in \omega$ be such that $f(x) \leq g(x)$ for all $x \geq k$. To determine whether $n \in 0''$ for $n \geq k$, ask $0'$ whether W_n gets an element after stage $g(n)$. The answer to this question is no if and only if $n \in 0''$. \square

Lemma 10. *There is a computable function $f(n, s)$ which satisfies the following properties.*

- (1) $f(n) = \liminf_s f(n, s)$.
- (2) For every $k > f(n)$, there is a stage s_k such that for all $t \geq s_k$ either $f(n, t) = f(n)$ or $f(n, t) > k$.

Proof. The following procedure to define $f(n, s)$ is uniform in n and takes place in stages s . Fix n , let I_j , $j < 2^{n+1}$, list the subsets of $[0, n]$ and let $I^* = \{j \leq n \mid |W_j| = \infty\}$. We think of each I_j as a potential guess at I^* .

At stage s , consider each I_j . If a new element is enumerated into W_i at stage s and $i \in I_j$, then we declare $i \in I_j$ to be *checked in I_j* . (Once i is checked in I_j , it remains checked in I_j at future stages until I_j acts and we explicitly declare i to be *unchecked in I_j* .) If all of the elements of I_j are checked in I_j , then we declare I_j to be *ready*. Notice that the empty set is ready at every stage.

Let j denote the index such that I_j is ready and I_j has acted least recently. (This index is unique unless there is more than one j such that I_j is ready and I_j has never acted. In that case, let I_j be the least such index.) We let I_j *act* at stage s as follows. Define

$$G(n, s) = \max\{t < s \mid W_{n,t} \neq W_{n,t+1}\}$$

$$f(n, s) = \max\{G(k, s) \mid k \leq n \wedge k \notin I_j\}$$

Declare I_j *unready* and *uncheck* each $i \in I_j$. (If $i \in I_k$ is checked in I_k for $k \neq j$, then i remains checked in I_k .)

This completes the description of $f(n, s)$. It remains to verify Properties 1 and 2. Let $S^* = \{s \mid I_j = I^* \text{ acts at } s\}$. By the construction, S^* is infinite and

$$\lim_{s \in S^*} f(n, s) = f(n).$$

To verify $\liminf_s f(n, s) = \lim_{s \in S^*} f(n, s)$ and Property 2, we need to consider what happens when each $I_j \neq I^*$ acts. We split this analysis into two cases: when $\exists i(i \in I_j - I^*)$ and when $I_j \subseteq I^*$ but $\exists i(i \in I^* - I_j)$.

First, suppose $\exists i(i \in I_j - I^*)$. Since $i \notin I^*$, W_i is finite and hence $i \in I_j$ is only checked finitely often in I_j . Therefore, I_j only acts finitely often and has no effect on $f(n, s)$ for large enough values of s .

Second, suppose $I_j \subseteq I^*$ and $\exists i(i \in I^* - I_j)$. Let $S_j = \{s \mid I_j \text{ acts at } s\}$. Since $I_j \subseteq I^*$, S_j is infinite. Since $i \in I^*$, W_i is infinite and $\lim_s G(i, s) = \infty$. Since $i \notin I_j$, at each $s \in S_j$, $f(n, s) \geq G(i, s)$. Therefore, $\lim_{s \in S_j} f(n, s) = \infty$.

Properties 1 and 2 follow easily from these observations. \square

We need to add one further piece to our approximation of $0''$. (By appealing to the s-m-n theorem, we could have skipped this step, but it is perhaps easier to continue the theme of dominating functions.) We let $h(n)$ be the least stage s for which $K_s \upharpoonright n + 1 = K \upharpoonright n + 1$. (Here, $K = \{e \mid \varphi_e(e) \text{ halts}\}$ denotes the usual halting set for the partial computable functions.) Since $h(n)$ is a Δ_2^0 -function, it has a computable approximation $h(n, s)$ such that $\lim_s h(n, s) = h(n)$. Finally, let $a(n, s)$ be the computable function defined by $a(n, s) = \max\{f(n, s), h(n, s)\}$.

Lemma 11. *The computable function $a(n, s)$ satisfies the following properties.*

- (1) $a(n) = \liminf_s a(n, s)$ exists and for all n , $a(n) \geq f(n), h(n)$.
- (2) For every $k > a(n)$, there is a stage s_k such that for all $t \geq s_k$, either $a(n, t) = a(n)$ or $a(n, t) > k$.
- (3) For any function $b(n)$ dominating $a(n)$, $0'' \leq_T b$. In particular, $0'' \leq_T a$.

Proof. Properties 1 and 2 follow from Lemma 10 and the fact that $h(n) = \lim_s h(n, s)$. Property 3 follows from the fact that if $b(n)$ dominates $a(n)$, then $b(n)$ dominates both $h(n)$ and $f(n)$. The fact that $b(n)$ dominates $h(n)$ gives $0' \leq_T b$. Combining this fact with Lemma 9 gives $0'' \leq_T b$. \square

We use the function $a(n, s)$ to prove Theorem 1. We start by fixing a *particular* family \mathcal{F} .

Definition 12. From now on, we fix the family \mathcal{F} to be the set $\{X_s \mid s \in \omega\}$ where X_s is the set of all pairs $\langle n, m \rangle$ such that $n \leq s$ and

$$\exists t \geq s [a(n, t) = m \wedge \forall u (s \leq u \leq t \rightarrow a(n, t) \leq a(n, u))].$$

Let $I_{\mathcal{F}} = \{\langle n, m \rangle \mid a(n) = m\}$.

Lemma 13. *\mathcal{F} is an a.e. family of finite subsets of ω with infinite limit $I_{\mathcal{F}}$.*

Proof. Each X_s is finite because $\langle n, m \rangle \in X_s$ implies that $n \leq s$ and $m \leq a(n, s)$. Therefore, \mathcal{F} is a family of finite sets. Furthermore, $I_{\mathcal{F}}$ is clearly infinite. It remains to show that \mathcal{F} is an a.e. family with limit $I_{\mathcal{F}}$.

On the one hand, assume that $\langle n, m \rangle \in I_{\mathcal{F}}$. We show that $\langle n, m \rangle \in X_s$ for almost all s . By the definition of $I_{\mathcal{F}}$, $a(n) = m$. By Lemma 11, we can fix s_0 such that for all $s \geq s_0$, $a(n, s) \geq a(n)$. For every $s \geq s_0$, there is a $t \geq s$ such that $a(n, t) = a(n)$ and hence the pair $\langle n, a(n) \rangle$ is eventually enumerated into X_s for each $s \geq s_0$.

On the other hand, assume that $\langle n, m \rangle \notin I_{\mathcal{F}}$. We show that $\langle n, m \rangle \in X_s$ for only finitely many s . By the definition of $I_{\mathcal{F}}$, $a(n) \neq m$. Since $a(n) = \liminf_s a(n, s)$, this claim is clear in the case when $m < a(n)$. If $a(n) < m$, then by Lemma 11, there is a stage s_m such that for all $s \geq s_m$, either $a(n, s) = a(n)$ or $a(n, s) > m$. Therefore, for all $s \geq s_m$, $a(n, s) \neq m$. It follows that $\langle n, m \rangle \notin X_s$ for any $s \geq s_m$. \square

By Lemma 5, the family \mathcal{F} generates a theory $T_{\mathcal{F}}$ which is trivial and strongly minimal.

Lemma 14. *There is a c.e. Friedberg enumeration of \mathcal{F} .*

Proof. This lemma follows because $a(n, s)$ is a computable function and hence the definition of X_s is uniformly Σ_1^0 . \square

By Lemma 7, $T_{\mathcal{F}}$ has a computable prime model.

Lemma 15. *If \mathcal{A} is a nonprime model of $T_{\mathcal{F}}$, then $0''$ is computable from the atomic diagram of \mathcal{A} .*

Proof. By Lemma 8, we can enumerate $I_{\mathcal{F}}$ from the atomic diagram of \mathcal{A} and hence we can enumerate the graph of $a(n)$ from the atomic diagram of \mathcal{A} . Therefore, we can compute the function $a(n)$ from the diagram of \mathcal{A} and thus can compute $0''$ by Lemma 11. \square

This completes the proof of Theorem 1.

2. SOME REMARKS ON AXIOMATIZABILITY

We now conclude with some remarks on axiomatizability by formulas of various quantifier-complexity.

Definition 16. An \mathcal{L} -theory T has a *quantifier-depth 1 axiomatization* if T is equivalent to a set S of sentences, each of which is a Boolean combination of universal sentences.

Definition 17. An \mathcal{L} -theory T is *Skolemized* if for every \mathcal{L} -formula $\varphi(x, \bar{y})$ there are finitely many terms $\tau_1(\bar{y}), \dots, \tau_n(\bar{y})$ such that

$$T \models \forall \bar{y} [\exists x \varphi(x, \bar{y}) \rightarrow \bigvee_i \varphi(\tau_i(\bar{y}), \bar{y})]$$

The Tarski-Vaught criterion for being an elementary substructure, coupled with an easy compactness argument, entails that T is Skolemized if and only if every substructure of every model of T is an elementary substructure. Recalling that a theory T has a universal axiomatization if and only if the class of models of T is closed under substructures, it follows immediately that any Skolemized theory has a universal axiomatization, and the converse holds whenever the theory is model complete.

Lemma 18. *The following are equivalent for any structure \mathcal{M} such that $Th(\mathcal{M}_M)$ is model complete:*

- (1) $Th(\mathcal{M}_M)$ has a quantifier-depth 1 axiomatization;
- (2) $Th(\mathcal{M}_M)$ has a universal axiomatization;
- (3) $Th(\mathcal{M}_M)$ is Skolemized.

Proof. (2) \Rightarrow (1) is trivial and (3) \Leftrightarrow (2) follows from our comments above. Finally, (1) \Rightarrow (2) holds for any structure \mathcal{M} . Indeed, suppose (1) holds for \mathcal{M} and choose $\mathcal{M}' \models Th(\mathcal{M}_M)$ and any \mathcal{L}_M -substructure \mathcal{N} of \mathcal{M}' . We may assume that $M \subseteq N \subseteq M'$ and $\mathcal{M} \preceq \mathcal{M}'$. It suffices to show that $\mathcal{N} \models Th(\mathcal{M}_M)$. Let σ be any universal \mathcal{L}_M -sentence. Clearly, if $\mathcal{N} \models \sigma$, then $\mathcal{M} \models \sigma$ as well. Conversely, if $\mathcal{M} \models \sigma$, then $\mathcal{M}' \models \sigma$ by elementarity, hence $\mathcal{N} \models \sigma$ as well. Thus, \mathcal{M} and \mathcal{N} agree on all universal \mathcal{L}_M -sentences, hence on all Boolean combinations of universal \mathcal{L}_M -sentences, so $\mathcal{N} \models Th(\mathcal{M}_M)$ by (1). \square

By adding dummy variables as needed, it is easily verified that for any \mathcal{L} -structure \mathcal{M} , $Th(\mathcal{M}_M)$ is Skolemized if and only if for every \mathcal{L} -formula $\varphi(x, \bar{y})$ there are finitely many \mathcal{L} -terms $\tau_1(\bar{y}, \bar{z}_1) \dots, \tau_n(\bar{y}, \bar{z}_n)$ such that

$$Th(\mathcal{M}) \models \exists \bar{z}_1 \dots \exists \bar{z}_n \forall \bar{y} [\exists x \varphi(x, \bar{y}) \rightarrow \bigvee_i \varphi(\tau_i(\bar{y}, \bar{z}_i), \bar{y})]$$

Thus, the property of $\text{Th}(\mathcal{M}_M)$ being Skolemized is preserved under elementary equivalence.

Proposition 19. *The following are equivalent for any trivial, strongly minimal (complete) \mathcal{L} -theory T :*

- (1) $\text{Th}(\mathcal{M}_M)$ has a quantifier-depth 1 axiomatization for some $\mathcal{M} \models T$;
- (2) $\text{Th}(\mathcal{M}_M)$ has a quantifier-depth 1 axiomatization for all $\mathcal{M} \models T$;
- (3) For all $M \subseteq A \subseteq M'$ such that $\mathcal{M} \preceq \mathcal{M}'$ are models of T , and for every $b \in \text{acl}_{\mathcal{M}'}(A)$, there are an \mathcal{L} -term $\tau(\bar{y})$ and $\bar{a} \in A$ with $\mathcal{M}' \models \tau(\bar{a}) = b$.

Proof. Recall from Goncharov/Harizanov/Laskowski/Lempp/McCoy [GHLLM03] that $\text{Th}(\mathcal{M}_M)$ is model complete for every model \mathcal{M} of T . Combining this with the preservation of being Skolemized under elementary equivalence, (1) \Leftrightarrow (2) follows from the Lemma.

Now assume that (2) holds. Choose $\mathcal{M}, \mathcal{M}', A, b$ as in (3) and let \mathcal{N} be the smallest substructure of \mathcal{M}' containing $A \cup \{b\}$. Since $\text{Th}(\mathcal{M}_M)$ is model complete, Lemma 18 implies that $\text{Th}(\mathcal{M}_M)$ is Skolemized, hence $\mathcal{N} \preceq \mathcal{M}'$. Thus $\text{acl}_{\mathcal{M}'}(A) = \text{acl}_{\mathcal{N}}(A) = N$, so (3) is satisfied.

Finally, suppose (3) holds and fix $\mathcal{M} \models T$. We show that $\text{Th}(\mathcal{M}_M)$ has a universal axiomatization by choosing $\mathcal{M}' \succeq \mathcal{M}$ and N such that $M \subseteq N \subseteq M'$ as \mathcal{L} -structures. By (3), N is infinite and algebraically closed in the \mathcal{L}_M -structure \mathcal{M}'_M , which has a strongly minimal theory. Thus $\mathcal{N} \models \text{Th}(\mathcal{M}_M)$ as required. \square

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