

# MATH 3795

## Lecture 5. Solving Linear Systems 3

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### Goals

- ▶ Positive definite and definite matrices.
- ▶ Cholesky decomposition.
- ▶ LU-Decomposition of Tridiagonal Systems
- ▶ Applications.

# Symmetric Positive Definite Matrices

- ▶  $A \in \mathbb{R}^{n \times n}$  is called symmetric if  $A = A^T$ .
- ▶  $A \in \mathbb{R}^{n \times n}$  is called symmetric positive definite if  $A = A^T$  and  $v^T A v > 0$  for all  $v \in \mathbb{R}^n$ ,  $v \neq 0$ .

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- ▶ Can we use the structure of  $A$ , i.e. symmetry and the positive definiteness of  $A$  or the property that  $A$  is  $m$ -banded to compute LU-decomposition more efficiently?

# Symmetric Positive Definite Matrices

- ▶ **Fact 1:** If  $L$  is a unit lower triangular matrix and  $U$  is an upper triangular matrix such that  $A = LU$ , then  $L$  and  $U$  are unique, i.e., if there  $\tilde{L}$  is a unit lower triangular matrix and  $\tilde{U}$  is an upper triangular matrix such that  $A = \tilde{L}\tilde{U}$ , then  $L = \tilde{L}$  and  $U = \tilde{U}$ .

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- ▶ **Fact 2:** If  $A \in \mathbb{R}^{n \times n}$  is symmetric positive definite and  $A = LU$ , then the diagonal entries of  $U$  are positive.
- ▶ Any upper triangular matrix  $U$  we can write as  $U = D\tilde{U}$ , where  $\tilde{U}$  is unit upper triangular and

$$D = \text{diag}(u_{11}, \dots, u_{nn}).$$

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Thus,

$$A = LU = LD\tilde{U}.$$

On the other hand

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Using that  $A = A^T$  and LU decomposition is unique

$$A = LU = \tilde{U}^T DL^T = \tilde{U}^T (DL^T) = (\text{lower unit triangular}) \times (\text{upper triangular}).$$

Thus,

$$L = \tilde{U}^T \quad \text{and} \quad U = DL^T.$$

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- ▶ Above we showed that if  $A$  is a symmetric positive definite matrix, then

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- ▶ Recall that

$$D = \text{diag}(u_{11}, \dots, u_{nn})$$

has positive diagonal entries. So we can define

$$D^{1/2} = \text{diag}(\sqrt{u_{11}}, \dots, \sqrt{u_{nn}}).$$

Define  $R := D^{1/2}L^T$ , then

$$A = LDL^T = LD^{1/2}D^{1/2}L^T = R^T R : \quad \text{Cholesky-decomposition}$$

## Example

$$\begin{aligned} \underbrace{\begin{pmatrix} 4 & -4 & 8 \\ -4 & 8 & -4 \\ 8 & -4 & 29 \end{pmatrix}}_A &= \underbrace{\begin{pmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 2 & 1 & 1 \end{pmatrix}}_L \underbrace{\begin{pmatrix} 4 & -4 & 8 \\ 0 & 4 & 4 \\ 0 & 0 & 9 \end{pmatrix}}_U \\ &= \underbrace{\begin{pmatrix} 1 & 0 & 0 \\ -1 & 1 & 0 \\ 2 & 1 & 1 \end{pmatrix}}_L \underbrace{\begin{pmatrix} 4 & 0 & 0 \\ 0 & 4 & 0 \\ 0 & 0 & 9 \end{pmatrix}}_D \underbrace{\begin{pmatrix} 1 & -1 & 2 \\ 0 & 1 & 1 \\ 0 & 0 & 1 \end{pmatrix}}_{L^T} \\ &= \underbrace{\begin{pmatrix} 2 & 0 & 0 \\ -2 & 2 & 0 \\ 4 & 2 & 3 \end{pmatrix}}_{R^T} \underbrace{\begin{pmatrix} 2 & -2 & 4 \\ 0 & 2 & 2 \\ 0 & 0 & 3 \end{pmatrix}}_R \end{aligned}$$

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- ▶ A matrix  $A \in \mathbb{R}^{n \times n}$  is symmetric positive definite if and only if there exists an upper triangular matrix  $R$  with  $r_{ii} > 0$ ,  $i = 1, \dots, n$ , such that  $A = R^T R$ .

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- ▶ So far we know how to compute  $LDL^T$  and  $R^T R$  by first computing the LU-decomposition and then derive  $LDL^T$  or  $R^T R$  from it.  
No savings. Not very useful.

# Cholesky Decomposition

- ▶ If  $A$  is symmetric positive definite, then there exists an upper triangular matrix  $R$  with  $r_{ii} > 0$ ,  $i = 1, \dots, n$ , such that  $A = R^T R$ .
- ▶ From the matrix-matrix multiplication we have

$$a_{ij} = \sum_{k=1}^n r_{ki} r_{kj} = \sum_{k=1}^{\min\{i,j\}} r_{ki} r_{kj}.$$

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- ▶ Can easily derive the algorithm
  - ▶ Fix  $i = 1$  and let  $j = 1 : n$ ;  $a_{1j} = r_{11} r_{1j}$ . This implies  $r_{11} = \sqrt{a_{11}}$ ,  $j = 1$  and  $r_{1j} = a_{1j}/r_{11}$ ,  $j = 2 : n$ .
  - ▶ Fix  $i = 2$  and let  $j = 2 : n$ ;  $a_{2j} = r_{12} r_{1j} + r_{22} r_{2j}$ . This implies  $r_{22} = \sqrt{a_{22} - r_{12} r_{12}}$ ,  $j = 2$ , and  $r_{2j} = (a_{2j} - r_{12} r_{1j})/r_{22}$ ,  $j = 3 : n$ .
  - ▶ Continue this way to determine all rows of  $R$  successively.

# Cholesky Decomposition

Input:  $A \in \mathbb{R}^{n \times n}$ , symmetric positive definite.

Output:  $R \in \mathbb{R}^{n \times n}$

1. For  $i = 1 : n$
2. If  $a_{ii} - \sum_{k=1}^{i-1} r_{ik}^2 \leq 0$ , return with an error message:  
 $A$  is not positive definite.
3.  $r_{ii} = \left( a_{ii} - \sum_{k=1}^{i-1} r_{ik}^2 \right)^{1/2}$
4. For  $j = i + 1 : n$
5.  $r_{ij} = \frac{1}{r_{ii}} \left( a_{ij} - \sum_{k=1}^{i-1} r_{ik} r_{kj} \right)^{1/2}$
6. End
7. End

Total number of flops:  $n^3/3$  (compare to  $2n^3/3$  for LU).

Matlab command syntax:  $[R] = chol(A)$ .

# Tridiagonal Systems

- ▶ A matrix  $A \in \mathbb{R}^{n \times n}$  is called a tridiagonal matrix if

$$a_{ij} = 0 \quad \text{for all } i, j \quad \text{with } |i - j| > 1.$$



$$A = \begin{pmatrix} d_1 & e_1 & & & & \\ c_1 & d_2 & e_2 & & & \\ & \ddots & \ddots & \ddots & & \\ & & & c_{n-2} & d_{n-1} & e_{n-1} \\ & & & & c_{n-1} & d_n \end{pmatrix}$$

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- ▶ There are only about  $3n$  entries in matrix  $A$ .
- ▶ **Question:** How can we compute the LU-decomposition of  $A$  efficiently?

# LU-Decomposition of Tridiagonal Systems

$$\underbrace{\begin{pmatrix} 1 & & & & & \\ -\frac{c_1}{d_1} & 1 & & & & \\ 0 & 0 & \ddots & & & \\ & & \ddots & \ddots & & \\ & & & \ddots & 1 & \\ & & & & 0 & 1 \end{pmatrix}}_{=M_1}, \quad \underbrace{\begin{pmatrix} d_1 & e_1 & & & & \\ & \tilde{d}_2 & e_2 & & & \\ & c_2 & d_3 & e_3 & & \\ & & \ddots & \ddots & \ddots & \\ & & & c_{n-2} & d_{n-1} & e_{n-1} \\ & & & & c_{n-1} & d_n \end{pmatrix}}_{=M_1 A, \quad \text{where } \tilde{d}_2 = d_2 - \frac{c_1}{d_1} e_1}$$





# LU-Decomposition of Tridiagonal Systems

Input:  $A \in \mathbb{R}^{n \times n}$ , symmetric positive definite.

Output:  $R \in \mathbb{R}^{n \times n}$

1. For  $k = 1 : n - 1$
2. If  $d_k = 0$ , stop
3.  $c_k = c_k / d_k$
4.  $d_{k+1} = d_{k+1} - c_k e_k$
5. End

This algorithm requires about  $3n$  flops.

## Application: Boundary Value Problems

- ▶ We want to compute a function  $y(x)$  that satisfies the differential equation

$$\begin{aligned} -\alpha y''(x) + \beta y'(x) + \gamma y(x) &= f(x), & x \in (a, b), \\ y'(a) &= y_a, \\ y(b) &= y_b. \end{aligned}$$

The system is called a two-point boundary value problem BVP.

- ▶ The slope of the unknown function  $y$  is specified at the left end-point  $a$  and the value of the unknown function  $y$  is specified at the right end-point  $b$ . Other boundary conditions are possible and will be considered later.
- ▶ The function  $f$  and the scalars  $\alpha$ ,  $\beta$ ,  $\gamma$ ,  $y_a$ , and  $y_b$  are given. We assume that  $\alpha > 0$  and  $\gamma \geq 0$ .

## Application: Boundary Value Problems

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- ▶ The solution  $y$  of a BVP is rarely known analytically, especially if  $\alpha$ ,  $\beta$ , and  $\gamma$  are nonconstant. In most application, the solution has to be approximated by numerical methods. We use the finite difference method.
- ▶ We select a grid

$$a = x_0 < x_1 < \cdots < x_n = b$$

with equidistant points

$$x_i = a + i \frac{b-a}{n} = a + i \Delta x$$

where the mesh size

$$\Delta x = \frac{b-a}{n}.$$

## Application: Boundary Value Problems

- ▶ **Finite differences** At a point  $x$  we can approximate the derivative of a function  $g : \mathbb{R} \rightarrow \mathbb{R}$  by

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- ▶ If we apply the approximation (1) with  $h = \frac{\Delta x}{2}$  to  $g(x) = y'(x)$  we obtain

$$-y''(x_i) = -(y'(x_i))' \approx -\frac{y'(x_i + \Delta x/2) - y'(x_i - \Delta x/2)}{\Delta x}.$$

We approximate the derivatives  $y'(x_i \pm \Delta x/2)$  using (1) with  $h = \Delta x/2$  and  $g(x) = y(x)$ . This gives

## Application: Boundary Value Problems

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$$\begin{aligned} -y''(x) &\approx -\frac{\frac{y(x_{i+1})-y(x_i)}{\Delta x} - \frac{y(x_i)-y(x_{i-1}))}{\Delta x}}{\Delta x} \\ &= \frac{-y(x_{i+1}) + 2y(x_i) - y(x_{i-1}))}{(\Delta x)^2}. \end{aligned} \tag{2}$$

Similarly,

$$y'(x) \approx \frac{y(x_{i+1}) - y(x_{i-1}))}{2\Delta x}. \tag{3}$$

## Application: Boundary Value Problems

Thus in the BVP

$$\begin{aligned} -\alpha y''(x) + \beta y'(x) + \gamma y(x) &= f(x), & x \in (a, b), \\ y'(a) &= y_a, \\ y(b) &= y_b. \end{aligned}$$

we can approximate at points  $x_i$ ,  $i = 0, 1, \dots, n$  by

$$\frac{-(\alpha + \frac{\Delta x}{2}\beta)y_{i-1} + (2\alpha + (\Delta x)^2\gamma)y_i - (\alpha - \frac{\Delta x}{2}\beta)y_{i+1}}{(\Delta x)^2} = f(x_i) \quad (4)$$

for  $i = 0 : n - 1$ . The boundary conditions and lead to

$$\frac{y_1 - y_{-1}}{2\Delta x} = y_a, \quad (5)$$

i.e.,  $y_{-1} = y_1 - 2\Delta x y_a$ . Finally, the other boundary condition yields

$$y_n = y_b. \quad (6)$$



## Matlab's spdiags

Given three arrays of length  $n$ ,  $c = [c_1; \dots; c_n]$ ,  $d = [d_1; \dots; d_n]$ ;  
 $e = [e_1; \dots; e_n]$ .

The Matlab command

$$A = \text{spdiags}([c, d, e], -1 : 1, n, n);$$

generates in a sparse matrix format a matrix

$$A = \begin{pmatrix} d_1 & e_2 & & & \\ c_1 & d_2 & e_3 & & \\ & \ddots & \ddots & \ddots & \\ & & c_{n-2} & d_{n-1} & e_n \\ & & & c_{n-1} & d_n \end{pmatrix}$$

Example

```
>> c = [1; 2; 3]; d = [4; 5; 6]; e = [7; 8; 9];
```

```
>> A = spdiags([c, d, e], -1:1, 3,3);
```

```
>> full(A)
```

```
ans =
```

```
4 8 0
```

```
1 5 9
```

```
0 2 6
```

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- ▶ To find other variations of LU and Cholesky decompositions type *help lu* and *help chol*
- ▶ There are better algorithms available if matrix  $A$  has a special structure, like *tridiagonal*, *banded*, *etc.*